Package 'adagio'

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Description The R package 'adagio' will provide methods and algorithms for (discrete) optimization, e.g. knapsack and subset sum procedures, derivative-free Nelder-Mead and Hooke-Jeeves minimization, and some (evolutionary) global optimization functions.

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NeedsCompilation no

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assignment

Linear Sum Assignment Problem

Description

Linear (sum) assignment problem, or LSAP.

Usage

assignment(cmat, dir = "min")

Arguments

cmat	quadratic (numeric) matrix, the cost matrix.
dir	direction, can be "min" or "max".

Details

Solves the linear (sum) assignment problem for quadratic matrices. Uses the lp.assign function from the lpSolve package, that is it solves LSAP as a mixed integer linear programming problem.

Value

List with components perm, the permutation that defines the minimum solution, min, the minimum value, and err is always 0, i.e. not used at the moment.

Note

Slower than the Hungarian algorithm in package clue.

References

Burkard, R., M. Dell'Amico, and S. Martello (2009). Assignment Problems. Society for Industrial and Applied Mathematics (SIAM).

Martello, S., and P. Toth (1990). Knapsack Problems: Algorithms and Computer Implementations. John Wiley & Sons, Ltd.

bpp_approx

See Also

clue::solve_LSAP

Examples

```
## Example similar to clue::solve_LSAP
set.seed(8237)
x <- matrix(sample(1:100), nrow = 10)</pre>
y <- assignment(x)</pre>
# show permutation and check minimum sum
y$perm
                                 # 7 6 10 5 8 2 1 4 9 3
y$min
                                 # 173
z <- cbind(1:10, y$perm)</pre>
x[z]
                                 # 16 9 49 6 17 14 1 44 10 7
y$min == sum(x[z])
                                 # TRUE
## Not run:
## Example: minimize sum of distances of complex points
n <- 100
x <- rt(n, df=3) + 1i * rt(n, df=3)</pre>
y <- runif(n) + 1i * runif(n)</pre>
cmat <- round(outer(x, y, FUN = function(x,y) Mod(x - y)), 2)</pre>
system.time(T1 <- assignment(cmat))</pre>
                                         # elapsed: 0.003
T1$min / 100
                                          # 145.75
## Hungarian algorithm in package 'clue'
library("clue")
system.time(T2 <- solve_LSAP(cmat))</pre>
                                          # elapsed: 0.014
sum(cmat[cbind(1:n, T2)])
                                          # 145.75
## End(Not run)
```

bpp_approx

Approximate Bin Packing

Description

Solves the Bin Packing problem approximately.

Usage

```
bpp_approx(S, cap, method = c("firstfit", "bestfit", "worstfit"))
```

Arguments

S	vector of weights (or sizes) of items.
сар	same capacity for all the bins.
method	which approximate method to use.

Details

Solves approximately the Bin Packing problem for numeric weights and bins, all having the same volume.

Possible methods are "firstfit", "bestfit", and "worstfit". "firstfit" tries to place each item as early as possible, "bestfit" such that the remaining space in the bin is as small as possible, and "worstfit" such that the remaining space is as big as possible.

Best results are achieved with the "bestfit" method. "firstfit" may be a reasonable alternative. For smaller and medium-sized data the approximate results will come quite close to the exact solution, see the examples.

In general, the results are much better if the items in S are sorted decreasingly. If they are not, an immediate warning is issued.

Value

A list of the following components:

nbins	minimum number of bins.
xbins	index of the bin each item is assigned to.
sbins	sum of item sizes in each bin.
filled	total volume filled in the bins (as percentage).

Note

The Bin Packing problem can be solved as a Linear Program. The formulation is a bit tricky, and it turned out 'lpSolve' does not solve medium-sized problems in acceptable time. (Tests with 'Rglpk' will follow.)

Author(s)

Hans W. Borchers

References

Silvano Martello. "Bin packing problems". In: 23rd Belgian Mathematical Optimization Workshop, La-Roche-en-Ardennes 2019.

See Also

Function binpacking in package 'knapsack' (on R-Forge).

Examples

```
## (1)
S <- c(50, 3, 48, 53, 53, 4, 3, 41, 23, 20, 52, 49)
cap <- 100
bpp_approx(S, cap, method = "bestfit")
## exact -- $nbins 4, filled 99.75 %
## firstfit -- $nbins 6, filled 66.5 %</pre>
```

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change_making

```
## bestfit -- $nbins 5, filled 79.8 %
## ! when decreasingly sorted, 'bestfit' with nbins = 4
## (2)
S <- c(100,99,89,88,87,75,67,65,65,57,57,49,47,31,27,18,13,9,8,1)
cap <- 100
bpp_approx(S, cap, method = "firstfit")
# firstfit: 12 bins; exact: 12 bins
## Not run:
## (3)
S <- c(99,99,96,96,92,92,91,88,87,86,
        85,76,74,72,69,67,67,62,61,56,
        52,51,49,46,44,42,40,40,33,33,
        30, 30, 29, 28, 28, 27, 25, 24, 23, 22,
        21,20,17,14,13,11,10, 7, 7, 3)
cap <- 100
bpp_approx(S, cap)
# exact: 25; firstfit: 25; bestfit: 25 nbins
## (4)
# 20 no.s in 1..100, capacity 100
set.seed(7013)
S <- sample(1:100, 20, replace = TRUE)</pre>
cap <- 100
bpp_approx(sort(S, decreasing = TRUE), cap, method = "bestfit")
# exact: 12 bins; firstfit and bestfit: 13; worstfit: 14 bins
## End(Not run)
```

change_making

Description

Solves the Change Making problem as an integer linear program.

Change Making Problem

Usage

```
change_making(items, value)
```

Arguments

items	vector of integer numbers greater than zero.
value	integer number

Details

The Change Making problem attempts to find a minimal combination of items that sum up to a given value. If the items are distinct positive integers, the solution is unique.

If the problem is infeasible, i.e. there is no such combination, the returned count is 0.

The problem is treated as an Integer Linear Program (ILP) and solved with the 1p solver in 1pSolve.

Value

Returns a list with components count, the number of items used to sum up to the value, and solution, the number of items used per item.

References

See the Wikipedia article on the "change making problem".

See Also

setcover

Examples

```
items = c(2, 5, 10, 50, 100)
value = 999
change_making(items, value)
## Not run:
solutions <- numeric(20)
for (m in 1:20) {
    sol <- change_making(items, m)
    solutions[m] <- sol$count
}
solutions
#> [1] 0 1 0 2 1 3 2 4 3 1 4 2 5 3 2 4 3 5 4 2
## End(Not run)
```

CMAES

Covariance Matrix Adaptation Evolution Strategy

Description

The CMA-ES (Covariance Matrix Adaptation Evolution Strategy) is an evolutionary algorithm for difficult non-linear non-convex optimization problems in continuous domain. The CMA-ES is typically applied to unconstrained or bounded constraint optimization problems, and search space dimensions between three and fifty.

CMAES

Usage

Arguments

par	objective variables initial point.
fun	objective/target/fitness function.
lower,upper	lower and upper bounds for the parameters.
sigma	coordinate wise standard deviation (step size).
stopfitness	stop if fitness < stopfitness (minimization).
stopeval	stop after stopeval number of function evaluations
	additional parameters to be passed to the function.

Details

The CMA-ES implements a stochastic variable-metric method. In the very particular case of a convex-quadratic objective function the covariance matrix adapts to the inverse of the Hessian matrix, up to a scalar factor and small random fluctuations. The update equations for mean and covariance matrix maximize a likelihood while resembling an expectation-maximization algorithm.

Value

Returns a list with components xmin and fmin.

Be patient; for difficult problems or high dimensions the function may run for several minutes; avoid problem dimensions of 30 and more!

Note

There are other implementations of Hansen's CMAES in package 'cmaes' (simplified form) and in package 'parma' as cmaes() (extended form).

Author(s)

Copyright (c) 2003-2010 Nikolas Hansen for Matlab code PURECMAES; converted to R by Hans W Borchers. (Hansen's homepage: www.cmap.polytechnique.fr/~nikolaus.hansen/)

References

Hansen, N. (2011). The CMA Evolution Strategy: A Tutorial. https://arxiv.org/abs/1604.00772

Hansen, N., D.V. Arnold, and A. Auger (2013). Evolution Strategies. J. Kacprzyk and W. Pedrycz (Eds.). Handbook of Computational Intelligence, Springer-Verlag, 2015.

See Also

cmaes::cmaes, parma::cmaes

Examples

```
## Not run:
## Polynomial minimax approximation of data points
## (see the Remez algorithm)
n <- 10; m <- 101
                           # polynomial of degree 10; no. of data points
xi <- seq(-1, 1, length = m)
yi <- 1 / (1 + (5*xi)^2)
                           # Runge's function
pval <- function(p, x)</pre>
                         # Horner scheme
    outer(x, (length(p) - 1):0, "^") %*% p
pfit <- function(x, y, n) # polynomial fitting of degree n</pre>
    qr.solve(outer(x, seq(n, 0), "^"), y)
fn1 <- function(p)</pre>
                            # objective function
    max(abs(pval(p, xi) - yi))
pf <- pfit(xi, yi, 10)
                       # start with a least-squares fitting
sol1 <- pureCMAES(pf, fn1, rep(-200, 11), rep(200, 11))</pre>
zapsmall(sol1$xmin)
# [1] -50.24826
                   0.00000 135.85352
                                          0.00000 -134.20107
                                                               0.00000
# [7]
      59.19315
                   0.00000 -11.55888
                                         0.00000
                                                    0.93453
print(sol1$fmin, digits = 10)
# [1] 0.06546780411
## Polynomial fitting in the L1 norm
## (or use LP or IRLS approaches)
fn2 <- function(p)</pre>
    sum(abs(pval(p, xi) - yi))
sol2 <- pureCMAES(pf, fn2, rep(-100, 11), rep(100, 11))</pre>
zapsmall(sol2$xmin)
# [1] -21.93238 0.00000 62.91083 0.00000 -67.84847
                                                          0.00000
# [7] 34.14398 0.00000 -8.11899 0.00000 0.84533
print(sol2$fmin, digits = 10)
# [1] 3.061810639
## End(Not run)
```

fminviz, flineviz Visualize Function Minimum

Description

Visualizes multivariate functions around a point or along a line between two points in R^n.

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hamiltonian

Usage

flineviz(fn, x1, x2, npoints = 51, scaled = 0.1)

Arguments

fn	multivariate function to be visualized.
x0,x1,x2	points in n-dimensional space.
nlines	number of lines to plot.
npoints	number of points used to plot a line.
scaled	scale factor to extend the line(s).

Details

fminviz vizualizes the behavior of a multivariate function fn around a point x0. It randomly selects nlines lines through x0 in Rⁿ and draws the curves of the function along these lines in one graph.

Curves that have at least one point below fn(x0) are drawn in red, all others in blue. The scale on the x-axis is the Euclidean distance in R^n. The scale factor can change it.

flineviz vizualizes the behavior of a multivariate function fn along the straight line between the points x1 and x2. Points x1 and x2 are also plotted.

Value

Plots a line graph and returns NULL (invisibly).

Examples

```
## Not run:
f1 <- function(x) x[1]^2 - x[2]^2
fminviz(f1, c(0, 0), nlines = 10)
f2 <- function(x) (1 - x[1])^2 + 100*(x[2] - x[1]^2)^2
flineviz(f2, c(0, 0), c(1, 1))
## End(Not run)
```

hamiltonian

Finds a Hamiltonian path or cycle

Description

A Hamiltionian path or cycle (a.k.a. Hamiltonian circuit) is a path through a graph that visits each vertex exactly once, resp. a closed path through the graph.

hamiltonian

Usage

hamiltonian(edges, start = 1, cycle = TRUE)

Arguments

edges	an edge list describing an undirected graph.
start	vertex number to start the path or cycle.
cycle	Boolean, should a path or a full cycle be found.

Details

hamiltonian() applies a backtracking algorithm that is relatively efficient for graphs of up to 30–40 vertices. The edge list is first transformed to a list where the i-th component contains the list of all vertices connected to vertex i.

The edge list must be of the form c(v1, v2, v3, v2, ...) meaning that there are edges $v1 \rightarrow v2$, $v3 \rightarrow v4$, etc., connecting these vertices. Therefore, an edge list has an even number of entries.

If the function returns NULL, there is no Hamiltonian path or cycle. The function does not check if the graph is connected or not. And if cycle = TRUE is used, then there also exists an edge from the last to the first entry in the resulting path.

If a Hamiltonian cycle exists in the graph it will be found whatever the starting vertex was. For a Hamiltonian path this is different and a successful search may very well depend on the start.

Value

Returns a vector containing vertex number of a valid path or cycle, or NULL if no path or cycle has been found (i.e., does not exist); If a cycle was requested, there exists an edge from the last to the first vertex in this list of edges.

Note

See the igraph package for more information about handling graphs and defining them through edge lists or other constructs.

Author(s)

Hans W. Borchers

References

Papadimitriou, Ch. H., and K. Steiglitz (1998). Optimization Problems: Algorithms and Complexity. Prentice-Hall/Dover Publications.

See Also

Package igraph

hamiltonian

Examples

```
## Dodekaeder graph
D20_edges <- c(
     1, 2, 1, 5, 1, 6, 2, 3, 2, 8, 3, 4, 3, 10, 4, 5, 4, 12,
     5, 14, 6, 7, 6, 15, 7, 8, 7, 16, 8, 9, 9, 10, 9, 17, 10, 11,
    11, 12, 11, 18, 12, 13, 13, 14, 13, 19, 14, 15, 15, 20, 16, 17, 16, 20,
    17, 18, 18, 19, 19, 20)
hamiltonian(D20_edges, cycle = TRUE)
# [1] 1 2 3 4 5 14 13 12 11 10 9 8 7 16 17 18 19 20 15 6
hamiltonian(D20_edges, cycle = FALSE)
# [1] 1 2 3 4 5 14 13 12 11 10 9 8 7 6 15 20 16 17 18 19
## Herschel graph
# The Herschel graph the smallest non-Hamiltonian polyhedral graph.
H11_edges <- c(
     1, 2, 1, 8, 1, 9, 1, 10, 2, 3, 2, 11, 3, 4, 3, 9, 4, 5,
     4, 11, 5, 6, 5, 9, 5, 10, 6, 7, 6, 11, 7, 8, 7, 10, 8, 11)
hamiltonian(H11_edges, cycle = FALSE)
# NULL
## Not run:
## Example: Graph constructed from squares
N <- 45 # 23, 32, 45
Q <- (2:trunc(sqrt(2*N-1)))^2
sq_edges <- c()</pre>
for (i in 1:(N-1)) {
    for (j in (i+1):N) {
       if ((i+j)
           sq_edges <- c(sq_edges, i, j)</pre>
    }
}
require(igraph)
sq_graph <- make_graph(sq_edges, directed=FALSE)</pre>
plot(sq_graph)
if (N == 23) {
    # does not find a path with start=1 ...
    hamiltonian(sq_edges, start=18, cycle=FALSE)
    # hamiltonian(sq_edges)
                                               # NULL
} else if (N == 32) {
    # the first of these graphs that is Hamiltonian ...
    # hamiltonian(sq_edges, cycle=FALSE)
    hamiltonian(sq_edges)
} else if (N == 45) {
    # takes much too long ...
    # hamiltonian(sq_edges, cycle=FALSE)
    hamiltonian(sq_edges)
}
## End(Not run)
```

Historize

Description

Provides storage for function calls.

Usage

Historize(fun, len = 0, ...)

Arguments

fun	Function of one or several variables.
len	If > 0 , input values will be stored, too.
	Additional parameters to be handed to the function.

Details

Historize() adds storage to the function. If len=0 then only function values will be stored in a vector, if len>0 then variables will be stored in a vector, and function values will be stored in a matrix, one line for every function call.

If Fn = Historize(...) is the 'historized' function, then the storage can be retrieved with Fn(), and can be cleared with Fn(NULL).

Filling the storage will take extra time and can slow down the function calls. Especially also storing the variables used in the call (with len>0) will make it considerably slower.

Functions can have multivariate output; the user is asked to take care of handling the output vector or matrix correctly. The function may even require additional parameters.

Value

Returns a list, \$input the input variables as matrix, \$H2 the function values as vector, \$nvars the number of input variables of the function, and \$ncalls the number or recorded function calls.

Note

Can also be applied to functions that output a vector (same length for every call).

Author(s)

Hans W. Borchers

See Also

trace

hookejeeves

Examples

```
f \leq function(x) sum(x^2)
F <- Historize(f, len = 1)</pre>
c( F(c(1,1)), F(c(1,2)), F(c(2,1)), F(c(2,2)) )
#> [1] 2 5 5 8
F()
#> $input
        [,1] [,2]
#>
#> [1,] 1
               1
#> [2,]
           1
                2
#> [3,]
           2
                1
#> [4,]
           2
                2
#>
#> $values
#> [1] 2 5 5 8
#>
#> $nvars
#> [1] 2
#>
#> $ncalls
#> [1] 4
F(NULL)
            # reset memory
## Rastrigin under Differential Evolution
Fn <- Historize(fnRastrigin)</pre>
dm <- 10
lb <- rep(-5.2, dm); ub <- -lb
sol <- simpleDE(Fn, lower = lb, upper = ub)</pre>
fvalues <- Fn()$values</pre>
fvals <- cummin(fvalues)</pre>
## Not run:
plot(fvalues, type = 'p', col = 7, pch = '.', cex = 2,
     main = "Simple DE optimization", xlab = '', ylab = '')
lines(fvals, col = 2, lwd = 2)
legend(length(fvalues), max(fvalues),
       c("Intermediate values", "cummulated min"),
       xjust = 1, col = c(7, 2), lwd = 2)
grid()
## End(Not run)
```

hookejeeves

Hooke-Jeeves Minimization Method

Description

An implementation of the Hooke-Jeeves algorithm for derivative-free optimization.

Usage

```
hookejeeves(x0, f, lb = NULL, ub = NULL,
            tol = 1e-08,
            target = Inf, maxfeval = Inf, info = FALSE, ...)
```

Arguments

x0	starting vector.
f	nonlinear function to be minimized.
lb, ub	lower and upper bounds.
tol	relative tolerance, to be used as stopping rule.
target	iteration stops when this value is reached.
maxfeval	maximum number of allowed function evaluations.
info	logical, whether to print information during the main loop.
	additional arguments to be passed to the function.

Details

This method computes a new point using the values of f at suitable points along the orthogonal coordinate directions around the last point.

Value

List with following components:

xmin	minimum solution found so far.
fmin	value of f at minimum.
fcalls	number of function evaluations.
niter	number of iterations performed.

Note

Hooke-Jeeves is notorious for its number of function calls. Memoization is often suggested as a remedy.

For a similar implementation of Hooke-Jeeves see the 'dfoptim' package.

References

C.T. Kelley (1999), Iterative Methods for Optimization, SIAM.

Quarteroni, Sacco, and Saleri (2007), Numerical Mathematics, Springer-Verlag.

See Also

neldermead

knapsack

Examples

```
## Rosenbrock function
rosenbrock <- function(x) {</pre>
   n <- length(x)</pre>
   x1 <- x[2:n]
   x2 <- x[1:(n-1)]
   sum(100*(x1-x2^2)^2 + (1-x2)^2)
}
hookejeeves(c(0,0,0,0), rosenbrock)
# $xmin
# [1] 1.000000 1.000001 1.000002 1.000004
# $fmin
# [1] 4.774847e-12
# $fcalls
# [1] 2499
# $niter
#[1] 26
hookejeeves(rep(0,4), lb=rep(-1,4), ub=0.5, rosenbrock)
# $xmin
# [1] 0.5000000 0.26221320 0.07797602 0.00608027
# $fmin
# [1] 1.667875
# $fcalls
# [1] 571
# $niter
# [1] 26
```

knapsack

0-1 Knapsack Problem

Description

Solves the 0-1 (binary) single knapsack problem.

Usage

knapsack(w, p, cap)

Arguments

W	integer vector of weights.
р	integer vector of profits.
сар	maximal capacity of the knapsack, integer too.

Details

knapsack solves the 0-1, or: binary, single knapsack problem by using the dynamic programming approach. The problem can be formulated as:

Maximize sum(x*p) such that $sum(x*w) \le cap$, where x is a vector with x[i] == 0 or 1.

Knapsack procedures can even solve subset sum problems, see the examples 3 and 3' below.

Value

A list with components capacity, profit, and indices.

Author(s)

HwB email: <hwborchers@googlemail.com>

References

Papadimitriou, C. H., and K. Steiglitz (1998). Combinatorial Optimization: Algorithms and Complexity. Dover Publications 1982, 1998.

Horowitz, E., and S. Sahni (1978). Fundamentals of Computer Algorithms. Computer Science Press, Rockville, ML.

See Also

knapsack::knapsack

Examples

```
# Example 1
p <- c(15, 100, 90, 60, 40, 15, 10, 1)
w <- c( 2, 20, 20, 30, 40, 30, 60, 10)
cap <- 102
(is <- knapsack(w, p, cap))</pre>
# [1] 1 2 3 4 6 , capacity 102 and total profit 280
## Example 2
p <- c(70, 20, 39, 37, 7, 5, 10)
w <- c(31, 10, 20, 19, 4, 3, 6)
cap <- 50
(is <- knapsack(w, p, cap))</pre>
# [1] 1 4 , capacity 50 and total profit 107
## Not run:
## Example 3: subset sum
p \le seq(2, 44, by = 2)^2
w <- p
is <- knapsack(w, p, 2012)</pre>
p[is$indices] # 16 36 64 144 196 256 324 400 576
## Example 3': maximize number of items
#
   w \le seq(2, 44, by = 2)^2
```

maxempty

maxempty

Maximally Empty Rectangle Problem

Description

Find the largest/maximal empty rectangle, i.e. with largest area, not containing given points.

Usage

maxempty(x, y, ax = c(0, 1), ay = c(0, 1))

Arguments

х, у	coordinates of points to be avoided.
ax, ay	left and right resp. lower and upper constraints.

Details

Find the largest or maximal empty two-dimensional rectangle in a rectangular area. The edges of this rectangle have to be parallel to the edges of the enclosing rectangle (and parallel to the coordinate axes). 'Empty' means that none of the points given are contained in the interior of the found rectangle.

Value

List with area and rect the rectangle as a vector usable for the rect graphics function.

Note

The algorithm has a run-time of $O(n^2)$ while there are run-times of O(n*log(n)) reported in the literature, utilizing a more complex data structure. I don't know of any comparable algorithms for the largest empty circle problem.

Author(s)

HwB email: <hwborchers@googlemail.com>

References

B. Chazelle, R. L. Drysdale, and D. T. Lee (1986). Computing the Largest Empty Rectangle. SIAM Journal of Computing, Vol. 15(1), pp. 300–315.

A. Naamad, D. T. Lee, and W.-L. Hsu (1984). On the Maximum Empty Rectangle Problem. Discrete Applied Mathematics, Vol. 8, pp. 267–277.

See Also

Hmisc::largest.empty with a Fortran implementation of this code.

Examples

```
maxquad
```

The MAXQUAD Test Function

Description

Lemarechal's MAXQUAD optimization test function.

Usage

maxquad(n, m)

Arguments

n	number of variables of the generated test function.
m	number of functions to compete for the maximum.

Details

MAXQUAD actually is a family of minimax functions, parametrized by the number n of variables and the number m of functions whose maximum it is.

maxsub

Value

Returns a list with components fn the generated test function of n variables, and gr the corresponding (analytical) gradient function.

References

Kuntsevich, A., and F. Kappel (1997). SolvOpt – The Solver for Local Nonlinear Optimization Problems. Manual Version 1.1, Institute of Mathematics, University of Graz.

Lemarechal, C., and R. Mifflin, Eds. (1978). Nonsmooth Optimization. Pergamon Press, Oxford.

Shor, N. Z. (1985). Minimization Methods for Non-differentiable Functions. Series in Computational Mathematics, Springer-Verlag, Berlin.

Examples

```
# Test function of 5 variables, defined as maximum of 5 smooth functions
maxq <- maxquad(5, 5)
fnMaxquad <- maxq$fn
grMaxquad <- maxq$gr
# shor</pre>
```

```
maxsub
```

Maximal Sum Subarray

Description

Find a subarray with maximal positive sum.

Usage

maxsub(x, inds = TRUE)

maxsub2d(A)

Arguments

Х	numeric vector.
А	numeric matrix
inds	logical; shall the indices be returned?

Details

maxsub finds a contiguous subarray whose sum is maximally positive. This is sometimes called Kadane's algorithm. maxsub will use a very fast version with a running time of O(n) where n is the length of the input vector x.

maxsub2d finds a (contiguous) submatrix whose sum of elements is maximally positive. The approach taken here is to apply the one-dimensional routine to summed arrays between all rows of A. This has a run-time of $O(n^3)$, though a run-time of $O(n^2 \log n)$ seems possible see the reference below. maxsub2d can solve a 100-by-100 matrix in a few seconds – but beware of bigger ones.

Value

Either just a maximal sum, or a list this sum as component sum plus the start and end indices as a vector inds.

Note

In special cases, the matrix A may be sparse or (as in the example section) only have one nonzero element in each row and column. Expectation is that there may exists a more efficient (say $O(n^2)$) algorithm in these special cases.

Author(s)

HwB <hwborchers@googlemail.com>

References

Bentley, Jon (1986). "Programming Pearls", Column 7. Addison-Wesley Publ. Co., Reading, MA.

T. Takaoka (2002). Efficient Algorithms for the Maximum Subarray Problem by Distance Matrix Multiplication. The Australasian Theory Symposion, CATS 2002.

Examples

```
## Find a maximal sum subvector
set.seed(8237)
x <- rnorm(1e6)
system.time(res <- maxsub(x, inds = TRUE))</pre>
res
## Standard example: Find a maximal sum submatrix
A <- matrix(c(0,-2,-7,0, 9,2,-6,2, -4,1,-4,1, -1,8,0,2),
            nrow = 4, ncol = 4, byrow =TRUE)
maxsub2d(A)
# $sum: 15
# $inds: 2 4 1 2 , i.e., rows = 2..4, columns = 1..2
## Not run:
## Application to points in the unit square:
set.seed(723)
N \leq 50; w \leq rnorm(N)
x <- runif(N); y <- runif(N)</pre>
clr <- ifelse (w >= 0, "blue", "red")
plot(x, y, pch = 20, col = clr, xlim = c(0, 1), ylim = c(0, 1))
xs <- unique(sort(x)); ns <- length(xs)</pre>
X <- c(0, ((xs[1:(ns-1)] + xs[2:ns])/2), 1)
ys <- unique(sort(y)); ms <- length(ys)</pre>
Y <- c(0, ((ys[1:(ns-1)] + ys[2:ns])/2), 1)
abline(v = X, col = "gray")
abline(h = Y, col = "gray")
A \leq matrix(0, N, N)
```

mknapsack

```
xi <- findInterval(x, X); yi <- findInterval(y, Y)
for (i in 1:N) A[yi[i], xi[i]] <- w[i]
msr <- maxsub2d(A)
rect(X[msr$inds[3]], Y[msr$inds[1]], X[msr$inds[4]+1], Y[msr$inds[2]+1])
## End(Not run)</pre>
```

mknapsack

Multiple 0-1 Knapsack Problem

Description

Solves the 0-1 (binary) multiple knapsack problem.

Usage

mknapsack(w, p, cap)

Arguments

W	vector of (positive) weights.
р	vector of (positive) profits.
сар	vector of capacities of different knapsacks.

Details

Solves the 0-1 multiple knapsack problem for a set of profits and weights. A multiple 0-1 knapsack problem can be formulated as:

```
 \begin{array}{l} \mbox{maximize vstar} = p(1) * (x(1,1) + \ldots + x(m,1)) + \ldots + p(n) * (x(1,n) + \ldots + x(m,n)) \mbox{ subject to } w(1) * x(i,1) + \ldots + w(n) * x(i,n) <= \mbox{cap}(i) \mbox{ for } i=1, \ldots, m \ x(1,j) + \ldots + x(m,j) <= \mbox{1 for } j=1, \ldots, n \ x(i,j) = 0 \mbox{ or } 1 \mbox{ for } i=1, \ldots, m \ x(j,j) \ x(j,j) = 0 \mbox{ for } i=1, \ldots, m \ x(j,j) \
```

The multiple knapsack problem is reformulated as a linear program and solved with the help of package lpSolve.

This function can be used for the single knapsack problem as well, but the 'dynamic programming' version in the knapsack function is faster (but: allows only integer values).

The solution found is most often not unique and may not be the most compact one. In the future, we will attempt to 'compactify' through backtracking. The number of backtracks will be returned in list element bs.

Value

A list with components, ksack the knapsack numbers the items are assigned to, value the total value/profit of the solution found, and bs the number of backtracks used.

Note

Contrary to earlier versions, the sequence of profits and weights has been interchanged: first the weights, then profits.

The compiled version was transferred to the knapsack package on R-Forge (see project 'optimist').

References

Kellerer, H., U. Pferschy, and D. Pisinger (2004). Knapsack Problems. Springer-Verlag, Berlin Heidelberg.

Martello, S., and P. Toth (1990). Knapsack Problems: Algorithms and Computer Implementations. John Wiley & Sons, Ltd.

See Also

Other packages implementing knapsack routines.

Examples

```
## Example 1: single knapsack
w <- c( 2, 20, 20, 30, 40, 30, 60, 10)
p <- c(15, 100, 90, 60, 40, 15, 10, 1)
cap <- 102
(is <- mknapsack(w, p, cap))</pre>
which(is$ksack == 1)
# [1] 1 2 3 4 6 , capacity 102 and total profit 280
## Example 2: multiple knapsack
w <- c( 40, 60, 30, 40, 20, 5)
p <- c(110, 150, 70, 80, 30, 5)
cap <- c(85, 65)
is <- mknapsack(w, p, cap)</pre>
# kps 1: 1,4; kps 2: 2,6; value: 345
## Example 3: multiple knapsack
p <- c(78, 35, 89, 36, 94, 75, 74, 79, 80, 16)
w <- c(18, 9, 23, 20, 59, 61, 70, 75, 76, 30)
cap <- c(103, 156)
is <- mknapsack(w, p, cap)</pre>
# kps 1: 3,4,5; kps 2: 1,6,9; value: 452
## Not run:
# How to Cut Your Planks with R
# R-bloggers, Rasmus Baath, 2016-06-12
#
# This is application of multiple knapsacks to cutting planks into pieces.
planks_we_have <- c(120, 137, 220, 420, 480)
planks_we_want <- c(19, 19, 19, 19, 79, 79, 79, 103, 103,
                    103, 135, 135, 135, 135, 160)
s <- mknapsack(planks_we_want, planks_we_want + 1, planks_we_have)</pre>
s$ksack
```

```
## [1] 5 5 5 5 3 5 5 4 1 5 4 5 3 2 4
# Solution w/o backtracking
# bin 1 : 103
                                     | Rest: 17
# bin 2 : 135
                                     | Rest: 2
# bin 3 : 79 + 135
                                     | Rest:
                                              6
# bin 4 : 103 + 135 + 160
                                     | Rest: 22
# bin 5 : 4*19 + 2*79 + 103 + 135
                                     | Rest:
                                              8
#
# Solution with reversing the bins (bigger ones first)
# bin 1 : 103
                                     | Rest:
                                             4
# bin 2 : 2*19 +
                 79
                                     | Rest:
                                             20
# bin 3 : 79 + 135
                                     | Rest:
                                              6
# bin 4 : 2*19 + 79 + 135 + 160
                                     | Rest:
                                              8
# bin 5 : 2*103 + 2*135
                                     | Rest: 17
#
# Solution with backtracking (compactification)
# sol = c(1, 4, 4, 1, 1, 3, 4, 5, 5, 5, 5, 4, 2, 3, 4)
# bin 1 : 2*19 + 79
                                     | Rest: 3
                                             2
# bin 2 : 135
                                     | Rest:
# bin 3 : 79 + 135
                                     | Rest: 6
# bin 4 : 2*19 + 79 + 135 + 160
                                              8
                                     | Rest:
# bin 5 : 3*103 + 135
                                     | Rest: 36
```

End(Not run)

neldermead

Nelder-Mead Minimization Method

Description

An implementation of the Nelder-Mead algorithm for derivative-free optimization / function minimization.

Usage

Arguments

fn	nonlinear function to be minimized.
x0	starting point for the iteration.
adapt	logical; adapt to parameter dimension.

tol	terminating limit for the variance of function values; can be made *very* small, like tol=1e-50.
maxfeval	maximum number of function evaluations.
step	size and shape of initial simplex; relative magnitudes of its elements should reflect the units of the variables.
	additional arguments to be passed to the function.
lower, upper	lower and upper bounds.

Details

Also called a 'simplex' method for finding the local minimum of a function of several variables. The method is a pattern search that compares function values at the vertices of the simplex. The process generates a sequence of simplices with ever reducing sizes.

The simplex function minimisation procedure due to Nelder and Mead (1965), as implemented by O'Neill (1971), with subsequent comments by Chambers and Ertel 1974, Benyon 1976, and Hill 1978. For another elaborate implementation of Nelder-Mead in R based on Matlab code by Kelley see package 'dfoptim'.

eldermead can be used up to 20 dimensions (then 'tol' and 'maxfeval' need to be increased). With adapt=TRUE it applies adaptive coefficients for the simplicial search, depending on the problem dimension – see Fuchang and Lixing (2012). This approach especially reduces the number of function calls.

With upper and/or lower bounds, neldermeadb applies transfinite to define the function on all of Rⁿ and to retransform the solution to the bounded domain. Of course, if the optimum is near to the boundary, results will not be as accurate as when the minimum is in the interior.

Value

List with following components:

xmin	minimum solution found.
fmin	value of f at minimum.
fcount	number of iterations performed.
restarts	number of restarts.
errmess	error message

Note

Original FORTRAN77 version by R O'Neill; MATLAB version by John Burkardt under LGPL license. Re-implemented in R by Hans W. Borchers.

References

Nelder, J., and R. Mead (1965). A simplex method for function minimization. Computer Journal, Volume 7, pp. 308-313.

O'Neill, R. (1971). Algorithm AS 47: Function Minimization Using a Simplex Procedure. Applied Statistics, Volume 20(3), pp. 338-345.

neldermead

J. C. Lagarias et al. (1998). Convergence properties of the Nelder-Mead simplex method in low dimensions. SIAM Journal for Optimization, Vol. 9, No. 1, pp 112-147.

Fuchang Gao and Lixing Han (2012). Implementing the Nelder-Mead simplex algorithm with adaptive parameters. Computational Optimization and Applications, Vol. 51, No. 1, pp. 259-277.

See Also

hookejeeves

Examples

```
## Classical tests as in the article by Nelder and Mead
# Rosenbrock's parabolic valley
rpv <- function(x) 100*(x[2] - x[1]^2)^2 + (1 - x[1])^2
x0 <- c(-2, 1)
                                        # 11
neldermead(rpv, x0)
# Fletcher and Powell's helic valley
fphv <- function(x)
    100*(x[3] - 10*atan2(x[2], x[1])/(2*pi))^2 +
        (sqrt(x[1]^2 + x[2]^2) - 1)^2 + x[3]^2
x0 <- c(-1, 0, 0)
neldermead(fphv, x0)
                                        # 100
# Powell's Singular Function (PSF)
psf <- function(x) (x[1] + 10*x[2])^2 + 5*(x[3] - x[4])^2 +
                    (x[2] - 2*x[3])^4 + 10*(x[1] - x[4])^4
x0 < -c(3, -1, 0, 1)
                            # 0 0 0 0, needs maximum number of function calls
neldermead(psf, x0)
# Bounded version of Nelder-Mead
lower <- c(-Inf, 0, 0)
upper <- c( Inf, 0.5, 1)
x0 <- c(0, 0.1, 0.1)
neldermeadb(fnRosenbrock, c(0, 0.1, 0.1), lower = lower, upper = upper)
# $xmin = c(0.7085595, 0.5000000, 0.2500000)
# $fmin = 0.3353605
## Not run:
# Can run Rosenbrock's function in 30 dimensions in one and a half minutes:
neldermead(fnRosenbrock, rep(0, 30), tol=1e-20, maxfeval=10^7)
# $xmin
# [1] 0.9999998 1.0000004 1.0000000 1.0000001 1.0000000 1.0000001
# [7] 1.0000002 1.0000001 0.9999997 0.9999999 0.9999997 1.0000000
# [13] 0.9999999 0.9999994 0.9999998 0.9999999 0.9999999 0.9999999
# [19] 0.9999999 1.0000001 0.9999998 1.0000000 1.0000003 0.9999999
# [25] 1.0000000 0.9999996 0.9999995 0.9999990 0.9999973 0.9999947
# $fmin
# [1] 5.617352e-10
# $fcount
# [1] 1426085
# elapsed time is 96.008000 seconds
```

occurs

End(Not run)

occurs

Description

Counts items, or finds subsequences of (integer) sequences.

Usage

```
count(x, sorted = TRUE)
```

occurs(subseq, series)

Arguments

х	array of items, i.e. numbers or characters.
sorted	logical; default is to sort items beforehand.
subseq	vector of integers.
series	vector of integers.

Details

count counts the items, similar to table, but as fast and a more tractable output. If sorted then the total number per item will be counted, else per repetition.

If m and n are the lengths of s and S resp., then occurs(s, S) determines all positions i such that s == S[i, ..., i+m-1].

The code is vectorized and relatively fast. It is intended to complement this with an implementation of Rabin-Karp, and possibly Knuth-Morris-Pratt and Boyer-Moore algorithms.

Value

count returns a list with components v the items and e the number of times it apears in the array. occurs returns a vector of indices, the positions where the subsequence appears in the series.

Examples

```
## Examples
patrn <- c(1,2,3,4)
exmpl <- c(3,3,4,2,3,1,2,3,4,8,8,23,1,2,3,4,4,34,4,3,2,1,1,2,3,4)
occurs(patrn, exmpl)
## [1] 6 13 23
## Not run:
set.seed(2437)
p <- sample(1:20, 1000000, replace=TRUE)</pre>
```

setcover

```
system.time(i <- occurs(c(1,2,3,4,5), p)) #=> [1] 799536
## user system elapsed
## 0.017 0.000 0.017 [sec]
system.time(c <- count(p))
## user system elapsed
## 0.075 0.000 0.076
print(c)
## $v
## [1] 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20
## $e
## [1] 49904 50216 49913 50154 49967 50045 49747 49883 49851 49893
## [11] 50193 50024 49946 49828 50319 50279 50019 49990 49839 49990</pre>
```

```
## End(Not run)
```

setcover

Set cover problem

Description

Solves the Set Cover problem as an integer linear program.

Usage

```
setcover(Sets, weights)
```

Arguments

Sets	matrix of 0s and 1s, each line defining a subset.
weights	numerical weights for each subset.

Details

The Set Cover problems attempts to find in subsets (of a 'universe') a minimal set of subsets that still covers the whole set.

Each line of the matrix Sets defines a characteristic function of a subset. It is required that each element of the universe is contained in at least one of these subsets.

The problem is treated as an Integer Linear Program (ILP) and solved with the lp solver in lpSolve.

Value

Returns a list with components sets, giving the indices of subsets, and objective, the sum of weights of subsets present in the solution.

References

See the Wikipedia article on the "set cover problem".

See Also

knapsack

Examples

SIAM test functions Trefethen and Wagon Test Functions

Description

Test functions for global optimization posed for the SIAM 100-digit challenge in 2002 by Nick Trefethen, Oxford University, UK.

Usage

fnTrefethen(p2)
fnWagon(p3)

Arguments

p2	Numerical vector of length 2.
р3	Numerical vector of length 3.

Details

These are highly nonlinear and oscillating functions in two and three dimensions with thousands of local mimima inside the unit square resp. cube (i.e., $[-1, 1] \times [-1, 1] \times [-1, 1] \times [-1, 1]$).

Value

Function value is a single real number.

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simpleDE

Author(s)

HwB <hwborchers@googlemail.com>

References

F. Bornemann, D. Laurie, S. Wagon, and J. Waldvogel (2004). The SIAM 100-Digit Challenge: A Study in High-Accuracy Numerical Computing. Society for Industrial and Applied Mathematics.

Examples

```
x <- 2*runif(5) - 1
fnTrefethen(x)
fnWagon(x)

## Not run:
T <- matrix(NA, nrow=1001, ncol=1001)
for (i in 1:1001) {
    for (j in 1:1001) {
        T[i, j] <- fnTrefethen(c(x[i], y[j]))
      }
    }
    image(x, y, T)
    contour(x, y, T, add=TRUE)

## End(Not run)</pre>
```

· ·

simpleDE

Simple Differential Evolution Algorithm

Description

Simple Differential Evolution for Minimization.

Usage

Arguments

fun	the objective function to be minimized.
lower	vector of lower bounds for all coordinates.
upper	vector of upper bounds for all coordinates.
Ν	population size.
nmax	bound on the number of generations.
r	amplification factor.
confined	logical; stay confined within bounds.
log	logical; shall a trace be printed.

Details

Evolutionary search to minimize a function: For points in the current generation, children are formed by taking a linear combination of parents, i.e., each member of the next generation has the form

$$p_1 + r(p_2 - p_3)$$

where the p_i are members of the current generation and r is an amplification factor.

Value

List with the following components:

fmin	function value at the minimum found.
xmin	numeric vector representing the minimum.
nfeval	number of function calls.

Note

Original Mathematica version by Dirk Laurie in the SIAM textbook. Translated to R by Hans W Borchers.

Author(s)

HwB <hwborchers@googlemail.com>

References

Dirk Laurie. "A Complex Optimization". Chapter 5 In: F. Bornemann, D. Laurie, S. Wagon, and J. Waldvogel (Eds.). The SIAM 100-Digit Challenge. Society of Industrial and Applied Mathematics, 2004.

See Also

simpleEA, DEoptim in the 'DEoptim' package.

Examples

```
simpleDE(fnTrefethen, lower = c(-1,-1), upper = c(1,1))
# $fmin
# [1] -3.306869
# $xmin
# [1] -0.02440308 0.21061243 # this is the true global optimum!
```

simpleEA

Description

Simple Evolutionary Algorithm for Minimization.

Usage

Arguments

fn	the objective function to be minimized.
lower	vector of lower bounds for all coordinates.
upper	vector of upper bounds for all coordinates.
Ν	number of children per parent.
	additional parameters to be passed to the function.
con	percentage of individuals concentrating to the best parents.
new	percentage of new individuals not focussed on existing parents.
tol	tolerance; if in the last three loops no better individuals were found up to this tolerance, stop.
eps	grid size bound to be reached.
scl	scaling factor for shrinking the grid.
confined	logical; shall the set of individuals be strictly respect the boundary? Default: FALSE.
log	logical, should best solution found be printed per step.

Details

Evolutionary search to minimize a function: For each point in the current generation, n random points are introduced and the n best results of each generation (and its parents) are used to form the next generation.

The scale shrinks the generation of new points as the algorithm proceeds. It is possible for some children to lie outside the given rectangle, and therefore the final result may lie outside the unit rectangle well. (TO DO: Make this an option.)

subsetsum

Value

List with the following components:

par	numeric vector representing the minimum found.
val	function value at the minimum found.
fun.calls	number of function calls made.
rel.scl	last scaling factor indicating grid size in last step.
rel.tol	relative tolerance within the last three minima found.

Note

Original Mathematica Version by Stan Wagon in the SIAM textbook. Translated to R by Hans W Borchers.

Author(s)

HwB <hwborchers@googlemail.com>

References

Stan Wagon. "Think Globally, Act Locally". Chapter 4 In: F. Bornemann, D. Laurie, S. Wagon, and J. Waldvogel (Eds.). The SIAM 100-Digit Challenge. Society of Industrial and Applied Mathematics, 2004.

See Also

DEoptim in the 'DEoptim' package.

Examples

```
simpleEA(fnTrefethen, lower=c(-1,-1), upper=c(1,1), log=FALSE)
# $par
# [1] -0.02440310 0.21061243 # this is the true global optimum!
# $val
# [1] -3.306869
```

subsetsum Subset Sum Problem

Description

Subset sum routine for positive integers.

Usage

```
subsetsum(S, t, method = "greedy")
```

sss_test(S, t)

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subsetsum

Arguments

S	vector of positive integers.
t	target value, bigger than all items in S.
method	can be "greedy" or "dynamic", where "dynamic" stands for the dynamic pro- gramming approach.

Details

subsetsum is searching for a set of elements in S that sum up to t by continuously adding more elements of S.

It is not required that S is decreasingly sorted. But for reasons of efficiency and smaller execution times it is urgently recommended to sort the item set in decreasing order. See the examples to find out how to handle your data.

The first components will be preferred, i.e., if S is decreasing, the sum with larger elements will be found, if increasing, the sum with smaller elements. Because of timing considerations, the default is to sort decreasingly before processing.

The dynamic method may be faster for large sets, but will also require much more memory if the target value is large.

sss_test will find the biggest number below or equal to t that can be expressed as a sum of items in S. It will not return any indices. It can be quite fast, though it preprocesses the set S to be sorted decreasingly, too.

Value

List with the target value, if reached, and vector of indices of elements in S that sum up to t.

If no solution is found, the dynamic method will return indices for the largest value below the target, the greedy method witll return NULL.

sss_test will simply return maximum sum value found.

Note

A compiled version – and much faster, in Fortran – can be found in package 'knapsack' (R-Forge, project 'optimist') as subsetsum. A recursive version, returning *all* solutions, is much too slow in R, but is possible in Julia and can be asked from the author.

Author(s)

HwB email: <hwborchers@googlemail.com>

References

Horowitz, E., and S. Sahni (1978). Fundamentals of Computer Algorithms. Computer Science Press, Rockville, ML.

See Also

maxsub

Examples

```
t <- 5842
S <- c(267, 493, 869, 961, 1000, 1153, 1246, 1598, 1766, 1922)
# S is not decreasingly sorted, so ...
o <- order(S, decreasing = TRUE)</pre>
So <- S[o]
                                    # So is decreasingly sorted
sol <- subsetsum(So, t)</pre>
                                    # $inds: 2 4 6 7 8 w.r.t. So
is <- o[sol$inds]</pre>
                                    # is:
                                               97543 w.r.t. S
sum(S[is])
                                    # 5842
## Not run:
amount <- 4748652
products <-
c(30500,30500,30500,30500,42000,42000,42000,42000,
  42000, 42000, 42000, 42000, 42000, 42000, 71040, 90900,
  76950,35100,71190,53730,456000,70740,70740,533600,
  83800, 59500, 27465, 28000, 28000, 28000, 28000, 28000,
  26140,49600,77000,123289,27000,27000,27000,27000,
  27000,27000,80000,33000,33000,55000,77382,48048,
  51186,40000,35000,21716,63051,15025,15025,15025,
  15025,800000,1110000,59700,25908,829350,1198000,1031655)
# prepare set
prods <- products[products <= amount] # no elements > amount
prods <- sort(prods, decreasing=TRUE) # decreasing order</pre>
# now find one solution
system.time(is <- subsetsum(prods, amount))</pre>
# user system elapsed
# 0.030
         0.000 0.029
prods[is]
# [1] 70740
                         71190
                                76950
                                        77382
                                                 80000
                                                          83800
                70740
# [8]
         90900 456000 533600 829350 1110000 1198000
sum(prods[is]) == amount
# [1] TRUE
# Timings:
#
              unsorted
                         decr.sorted
# "greedy"
                22.930
                            0.030
                                       (therefore the default settings)
# "dynamic"
                 2.515
                               0.860
                                         (overhead for smaller sets)
# sss_test
                 8.450
                               0.040
                                        (no indices returned)
## End(Not run)
```

Testfunctions

Optimization Test Functions

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Testfunctions

Description

Simple and often used test function defined in higher dimensions and with analytical gradients, especially suited for performance tests. Analytical gradients, where existing, are provided with the gr prefix. The dimension is determined by the length of the input vector.

Usage

```
fnRosenbrock(x)
grRosenbrock(x)
fnRastrigin(x)
grRastrigin(x)
fnNesterov(x)
grNesterov(x)
fnNesterov1(x)
fnNesterov2(x)
fnHald(x)
grHald(x)
fnShor(x)
grShor(x)
```

Arguments

х

numeric vector of a certain length.

Details

Rosenbrock – Rosenbrock's famous valley function from 1960. It can also be regarded as a least-squares problem:

$$\sum_{i=1}^{n-1} (1-x_i)^2 + 100(x_{i+1}-x_i^2)^2$$

No. of Vars.:	n >= 2
Bounds:	-5.12 <= xi <= 5.12
Local minima:	at $f(-1, 1,, 1)$ for $n \ge 4$
Minimum:	0.0
Solution:	xi = 1, i = 1:n

Nesterov – Nesterov's smooth adaptation of Rosenbrock, based on the idea of Chebyshev polynomials. This function is even more difficult to optimize than Rosenbrock's:

$$(1-x_1)^2/4 + \sum_{i=1}^{n-1} (1+x_{i+1}-2x_i^2)^2$$

Two nonsmooth Nesterov functions are available: Nesterov2 and Nesterov1, defined as

$$(1-x_1)^2/4 + \sum_{i=1}^{n-1} |1+x_{i+1}-2x_i^2|$$

Testfunctions

$$|1 - x_1|/4 + \sum_{i=1}^{n-1} (|1 + x_{i+1} - 2|x_i|)$$

No. of Vars.:	n >= 2
Bounds:	-5.12 <= xi <= 5.12
Local minima: ?	
Minimum:	0.0
Solution:	xi = 1, i = 1:n

Nesterov1 and **Nesterov2** – Simlar to Nesterov, except the terms added are taken with absolute value, which makes this function nonsmooth and painful for gradient-based optimization routines; no gradient provided.

(Nesterov2 uses absolute instead of quadratic terms.)

Rastrigin – Rastrigin's function is a famous, non-convex example from 1989 for global optimization. It is a typical example of a multimodal function with many local minima:

$$10n + \sum_{1}^{n} (x_i^2 - 10\cos(2\pi x_i))$$

No. of Vars.: $n \ge 2$ Bounds: $-5.12 \le xi \le 5.12$ Local minima:manyMinimum:0.0Solution:xi = 0, i = 1:n

Hald – Hald's function is a typical example of a non-smooth test function, from Hald and Madsen in 1981.

$$\max_{1 \le i \le n} \left| \frac{x_1 + x_2 t_i}{1 + x_3 t_i + x_4 t_i^2 + x_5 t_i^3} - \exp(t_i) \right|$$

where n = 21 and $t_i = -1 + (i - 1)/10$ for $1 \le i \le 21$.

 No. of Vars.:
 n =5

 Bounds:
 -1 <= xi <= 1</td>

 Local minima:
 ?

 Minimum:
 0.0001223713

 Solution:
 (0.99987763, 0.25358844, -0.74660757, 0.24520150, -0.03749029)

Shor – Shor's function is another typical example of a non-smooth test function, a benchmark for Shor's R-algorithm.

transfinite

Value

Returns the values of the test function resp. its gradient at that point. If an analytical gradient is not available, a function computing the gradient numerically will be provided.

References

Search the Internet.

Examples

```
x <- runif(5)
fnHald(x); grHald(x)
# Compare analytical and numerical gradient
shor_gr <- function(x) adagio:::ns.grad(fnShor, x) # internal gradient
grShor(x); shor_gr(x)</pre>
```

transfinite

Boxed Region Transformation

Description

Transformation of a box/bound constrained region to an unconstrained one.

Usage

```
transfinite(lower, upper, n = length(lower))
```

Arguments

lower, upper	lower and upper box/bound constraints.
n	length of upper, lower if both are scalars, to which they get repeated.

Details

Transforms a constraint region in n-dimensional space bijectively to the unconstrained R^n space, applying a atanh resp. exp transformation to each single variable that is bound constraint.

It provides two functions, h: $B = []x \dots x[] \longrightarrow R^n$ and its inverse hinv. These functions can, for example, be used to add box/bound constraints to a constrained optimization problem that is to be solved with a (nonlinear) solver not allowing constraints.

Value

Returns to functions as components h and hinv of a list.

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Based on an idea of Ravi Varadhan, intrinsically used in his implementation of Nelder-Mead in the 'dfoptim' package.

For positivity constraints, $x \ge 0$, this approach is considered to be numerically more stable than $x - \ge \exp(x)$ or $x - \ge x^2$.

Examples

```
lower <- c(-Inf, 0, 0)
upper <- c( Inf, 0.5, 1)
Tf <- transfinite(lower, upper)</pre>
h <- Tf$h; hinv <- Tf$hinv</pre>
## Not run:
## Solve Rosenbrock with one variable restricted
rosen <- function(x) {</pre>
    n <- length(x)
    x1 <- x[2:n]; x2 <- x[1:(n-1)]
    sum(100*(x1-x2^2)^2 + (1-x2)^2)
}
f <- function(x) rosen(hinv(x))</pre>
                                     # f must be defined on all of R^n
x0 <- c(0.1, 0.1, 0.1)
                                     # starting point not on the boundary!
nm <- nelder_mead(h(x0), f)</pre>
                                     # unconstraint Nelder-Mead
hinv(nm$xmin); nm$fmin
                                     # box/bound constraint solution
# [1] 0.7085596 0.5000000 0.2500004
# [1] 0.3353605
```

End(Not run)

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